Principles and Practice of Application Performance Measurement and Analysis on Parallel Systems

Lecture 1: Terminology and Methodology

Performance Tuning: an Old Problem!

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Performance Tuning: an Even Older Problem!!

"The most constant difficulty in contriving the engine has arisen from the desire to reduce the time in which the calculations were executed to the shortest which is possible."

Charles Babbage 1791 - 1871

Motivation

- High complexity in parallel and distributed systems
  - Four layers
    - Application
      - Algorithm, data structures
    - Parallel programming interface / Middleware
      - Compiler, parallel libraries, communication, synchronization
    - Operating system
      - Process and memory management, IO
    - Hardware
      - CPU, memory, network
  - Mapping/interaction between different layers
Performance Properties of Parallel Programs

Factors which influence performance of parallel programs

- “Sequential” factors
  - Computation
    - Choose right algorithm, use optimizing compiler
  - Cache and memory
    - Tough! Not many tools yet, hope compiler gets it right
  - Input / output
    - Not given enough attention

- “Parallel” factors
  - Communication (Message passing)
  - Threading
  - Synchronization
    - More or less understood, tool support

Performance Measurement Cycle

- Instrumentation
  - Insertion of extra code (probes, hooks) into application

- Measurement
  - Collection of data relevant to performance analysis

- Analysis
  - Calculation of metrics, identification of performance problems

- Presentation
  - Transformation of the results into a representation that can be easily understood by a human user

- Optimization
  - Elimination of performance problems
CONTENT

- Metrics
- Instrumentation techniques
  - Source code instrumentation
  - Binary instrumentation
- Instrumentation of parallel programs
  - MPI
  - OpenMP
- Measurement techniques
  - Profiling
  - Tracing

Metrics of Performance

- What can be measured?
  - A count of how many times an event occurs
    - E.g., Number of input / output requests
  - The duration of some time interval
    - E.g., duration of these requests
  - The size of some parameter
    - Number of bytes transmitted or stored

- Derived metrics
  - E.g., rates / throughput
  - Needed for normalization
Example Metrics

- Clock rate
- MIPS
  - Millions of instructions executed per second
- FLOPS
  - Floating-point operations per second
- Benchmarks
  - Standard test program(s)
  - Standardized methodology
  - E.g., SPEC, Linpack
- QUIPS / HINT [Gustafson and Snell, 95]
  - Quality improvements per second
  - Quality of solution instead of effort to reach it
- Execution time

Execution Time

- Wall-clock time
  - Includes waiting time: IO, memory, other system activities
  - In time-sharing environments also time consumed by other applications
- CPU time
  - Time spent by the CPU to execute the program
  - Execution time on behalf of the program
  - Does not include time the program was context-switched out
    - Problem: does not include inherent waiting time (e.g., IO)
    - Problem: portability? What is user, what is system time?
- Problem: execution time is non-deterministic
  - Use mean or minimum of several runs
Load Imbalance Metrics

- Imbalance Time
  - Metric time to identify code regions that need optimization
  - Two variations:
    - Computation Imbalance Time
      \[ \text{Computation Imbalance Time} = \text{Max Time} - \text{Avg time} \]
    - Synchronization Imbalance Time
      \[ \text{Synchronization Imbalance Time} = \text{Avg Time} - \text{Min time} \]
  - Provides an estimation to the user of how much time in the overall program would be saved if the corresponding section of the code had a perfect balance
  - Represents an upper bound on the “potential saving”

Load Imbalance Metrics

- Imbalance %
  - Provide an idea of the “badness” of the imbalance
  - Corresponds to the % of the time that the rest of the team, excluding the slowest PE is not engaged in useful work on the given function
  - “Percentage of resources available for parallelism” that is wasted

\[
\text{Imbalance\%} = 100 \times \frac{\text{Imbalance time}}{\text{Max Time}} \times \frac{N}{N - 1}
\]
**Speedup and Efficiency**

For a given problem A, let

- $\text{SerTime}(n)$ = Time of the best serial algorithm to solve A for input of size $n$
- $\text{ParTime}(n,p)$ = Time of the parallel algorithm + architecture to solve A for input size $n$, using p processors
- Note that $\text{SerTime}(n) \leq \text{ParTime}(n,1)$

Then

- $\text{Speedup}(p) = \frac{\text{SerTime}(n)}{\text{ParTime}(n,p)}$
- $\text{Work}(p) = p \cdot \text{ParTime}(n,p)$
- $\text{Efficiency}(p) = \frac{\text{SerTime}(n)}{[p \cdot \text{ParTime}(n,p)]}$

**Speedup and Efficiency II**

In general, expect

- $0 \leq \text{Speedup}(p) \leq p$
- Serial work $\leq$ Parallel work $< \infty$
- $0 \leq \text{Efficiency} \leq 1$

**Linear speedup**: if there is a constant $c > 0$ so that speedup is at least $c \cdot p$. Many use this term to mean $c = 1$.

**Perfect or ideal speedup**: speedup$(p) = p$

**Superlinear speedup**: speedup$(p) > p$ (efficiency $> 1$)

- Typical reason: Parallel computer has $p$ times more memory (cache), so higher fraction of program data fits in memory instead of disk (cache instead of memory)
- Parallel version is solving slightly different, easier problem or provides slightly different answer
**Amdahl’s Law**

- **Amdahl [1967]** noted:
  - Given a program, let \( f \) be fraction of time spent on operations that must be performed serially (unparallelizable work). Then for \( p \) processors:
    \[
    \text{Speedup}(p) \leq \frac{1}{f + (1-f)/p}
    \]
  - Thus no matter how many processors are used
    \[
    \text{Speedup}(p) \leq 1/f
    \]
  - Unfortunately, typical \( f \) is 5 – 20%

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**Maximal Possible Speedup / Efficiency**

![Graph showing maximal possible speedup and efficiency](image)

- \( f=0.001 \)
- \( f=0.01 \)
- \( f=0.1 \)
Amdahl’s Law II

- **Amdahl was an optimist**
  - Parallelization might require extra work, typically
    - Communication
    - Synchronization
    - Load balancing
  - Amdahl convinced many people that general-purpose parallel computing was not viable

- **Amdahl was a pessimist**
  - Fortunately, we can break the law!
  - Find better (parallel) algorithms with much smaller values of f
  - Superlinear speedup because of more data fits cache/memory
  - Scaling: time spent in serial portion is often a decreasing fraction of the total time as problem size increase

Scaling

- Sometimes the serial portion
  - is a fixed amount of time independent of problem size
  - or grows with problem size but slower than total time
  - Thus can often exploit large parallel machines by scaling the problem size with the number of processes

- Scaling approaches used for speedup reporting/measurements:
  - Fixed problem size (\(\Rightarrow\) strong scaling)
  - Fixed problem size per processor (\(\Rightarrow\) weak scaling)
  - Fixed time, find largest problem solvable [Gustafson 1988]
    - Commonly used in evaluating databases (transactions/s)
  - Fixed efficiency: find smallest problem to achieve it (\(\Rightarrow\) isoefficiency analysis)
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  - Binary instrumentation
- Instrumentation of parallel programs
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  - OpenMP
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  - Profiling
  - Tracing

Performance Tools Challenge

- User’s mental model of the program does not match the executed version
- Performance tools must be able to revert this semantic gap
Semantic Gap

- **Instrumentation levels**
  - Source code
  - Library
  - Runtime system
  - Object code
  - Operating system
  - Runtime image
  - Virtual machine

- **Problem**
  - Every level provides different information
  - Often instrumentation on multiple levels required

- **Challenge**
  - Mapping performance data onto application-level abstraction

Instrumentation Techniques

- **Static instrumentation**
  - Program is instrumented **prior to execution**

- **Dynamic instrumentation**
  - Program is instrumented **at runtime**

- **Code is inserted**
  - Manually
  - Automatically
    - By preprocessor / source-to-source translation tool
    - By compiler
    - By linking against pre-instrumented library or runtime system
    - By binary-rewrite / dynamic instrumentation tool

  e.g., “printf” ⇒ manual static source-code instrumentation
Source Code Instrumentation (I)

- For large complex applications, manual instrumentation is too tedious and error-prone ⇒ Tool support needed
- Automatic performance Instrumentation typically requires full source code parsers, e.g.,
  - Fortran, C: find 1st executable line and all exit points
  - C: executable code inside return statements

```c
int func(...) {
    double d;
    return (foo() * bar());
}
```

Source Code Instrumentation (II)

- Example C++ issues:
  - Template instrumentation?
  - Executing code before main
  - Function overloading
  - Operator overloading
- C++ instrumentation trick
  - Define instrumentation object
    ```c
    class Tracer { public:
            Tracer(...) { trace_enter(); }
            ~Tracer() { trace_exit(); }
        };
    ```
  - Declare instrumentation object as 1st statement in every function and method to be instrumented
    ```c
    int func(...) { Tracer trc_1;
        double d;
        return (foo() * bar());
    }
    ```
TAU Source Code Instrumentor

- Part of the TAU performance framework
- Supports
  - f77, f90
  - C, and C++
- Inserts calls to the TAU monitoring API
- Based on the Program Database Toolkit

http://tau.uoregon.edu/

Program Database Toolkit

- Based on commercial parsers
  - C, C++: Edison Design Group (EDG)
    - Full ISO 1998 C++ and ISO 1999 C Support
  - Fortran 77, Fortran90: Mutek, [Cleanscape]

- Program Database Utilities and Conversion Tools APplication Environment (DUCTAPE)
- Object-oriented Access to Static Information
- Classes, Modules, Routines, Types, Templates, Files, Macros, Namespaces, Comments/Pragmas, Statements (C/C++ only)

http://www.cs.uoregon.edu/research/pdt/
PDT Architecture and Tools

Binary Instrumentation

- Static binary rewrite
  - Instrumentation code is inserted into the binary before it starts to execute
  - Creates modified executable

- Dynamic binary instrumentation
  - On-the-fly: Insert, remove, and change instrumentation in the application program while it is running
  - Most flexible (but most complex) technique
  - Parallel programs
    - Coordinated instrumentation of all images needed
Dyninst

- Dyninst is a C++ library for machine-independent
  - process control and manipulation
  - runtime code generation
  - and binary patching
- University of Wisconsin and University of Maryland
- Basis for Paraview, DPCL, and OpenSpeedShop
- Open source
- Supports
  - Power/PowerPC (Linux)
  - X86 (Linux, BSD, Windows)
  - BlueGene/P
  - X86_64 (Linux, BSD, Windows)

- http://www.dyninst.org

Comparison of Techniques (I)

- Source code instrumentation
  - Portable
  - Link back to source code easy
  - Only way to capture "high-level" user abstractions

  - Recompilation necessary for
    (change in) instrumentation
  - Requires source code to be available
  - Hard to use for mixed-language applications
  - Source-to-source translation tool hard to implement
    for C++ and Fortran90
Comparison of Techniques (II)

- **Binary code instrumentation**
  - 😊 / 😞 The other way around compared to source instrumentation

- **Pre-instrumented library / runtime**
  - 😊 Easy to use: only re-linking necessary
  - 😞 Can only record information about library / runtime entities

- No single technique is sufficient!
- Typically, combinations of techniques needed!

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  - MPI
  - OpenMP
- Measurement techniques
  - Profiling
  - Tracing
Instrumentation of Parallel Programs

- **User-level** constructs
  - Modules / components / …
  - Program phases
  - Functions
  - Loops
  - …

- Constructs of the **parallel programming models**
  - Message passing
    - MPI, PVM, …
  - Threading and synchronization
    - OpenMP, POSIX, Win32, or Java threads, …

Instrumentation of User Functions

- **Ideally**: instrumentation by compiler or tool
  - Hidden, unsupported compiler options
    (GNU, Intel, IBM, NEC, Sun Fortran, PGI, Hitachi, ???)
  - TAU Source Code Instrumentor
  - TAU Binary Instrumentor (Dyninst)
  - TAU Virtual Machine Instrumentor (Java, Python)

- **Always works**: manually
  - Instrumentation APIs of tools: Scalasca, Vampirtrace, TAU, …
  - Scalasca’s POMP Directives
  - More details later …

- **Main problem**: selection of relevant constructs
PMPI: The MPI Profiling Interface

- Every MPI function has two names: MPI_xxx and PMPI_xxx
- This allows selective replacement of MPI routines at link time 
  ⇒ no re-compilation necessary
- Also called: wrapper function library
- Used by basically every MPI performance tools
  - VampirTrace, MPICH MPE, Scalasca EPIK, TAU, …

![Diagram of PMPI usage](image)

PMPI Example (C/C++)

```c
#include <stdio.h>
#include "mpi.h"

static int numsend = 0;

int MPI_Send(void *buf, int count, MPI_Datatype type,
              int dest, int tag, MPI_Comm comm) {
    numsend++;
    return PMPI_Send(buf, count, type, dest, tag, comm);
}

int MPI_Finalize() {
    int me;
    PMPI_Comm_rank(MPI_COMM_WORLD, &me);
    printf("%d sent %d messages.\n", me, numsend);
    return PMPI_Finalize();
}
```
PMPI Wrapper Development

- MPI has many functions! [MPI-1: 130  MPI-2: 320]
  - use wrapper generator (e.g., from MPICH MPE)
  - needed for Fortran and C/C++

- Message analysis / recording
  - Location recording ⇒ use ranks in MPI_COMM_WORLD?
  - Data volume ⇒ #elements * sizeof(type)
  - No message ID ⇒ need complete recording of traffic
  - Wildcard source and tag ⇒ record real values
  - Collective communication ⇒ communicator tracking
  - Non-blocking, persistent communication ⇒ track requests
  - Non-blocking ⇒ record recv at Wait*, Test*, Irecv?
  - One-sided communication?

OpenMP Monitoring?

- Problem:
  - OpenMP does not define standard monitoring interface
  - OpenMP is defined mainly by directives/pragmas

- Solution:
  - POMP: OpenMP Monitoring Interface
  - Joint Development
    - Forschungszentrum Jülich
    - University of Oregon
  - Presented at EWOMP’01, LACSI’01 and SC’01

Example:

```c
!$OMP PARALLEL DO POMP Instrumentation

call pomp_parallel_fork(d1)
!$OMP PARALLEL other-clauses...
call pomp_parallel_begin(d1)
call pomp_do_enter(d2)
!$OMP DO schedule-clauses, ordered-clauses,
  lastprivate-clauses
do loop
!$OMP END DO NOWAIT
call pomp_barrier_enter(d3)
!$OMP BARRIER
call pomp_barrier_exit(d3)
call pomp_do_exit(d2)
call pomp_parallel_end(d1)
!$OMP END PARALLEL DO
call pomp_parallel_join(d1)
```

### POMP-like Hooks in Production Compilers

- POMP was the base for the OpenMP instrumentation hooks provided in production compilers
  - Cray Compiling Environment
  - PGI
  - IBM XL compilers

- These instrumentation hooks are used for performance analysis of OpenMP in production tools
  - CrayPat
  - PGProf

- Also: New OpenMP ARB sanctioned low-level tool interface
  - [http://www.compunity.org/futures/omp-api.html](http://www.compunity.org/futures/omp-api.html)
  - Proof-of-concept implementations by Sun and Intel compilers
POMP Instrumentation Tool

- **OpenMP Pragma And Region Instrumentor**
- Source-to-source translator to insert POMP calls around OpenMP constructs and API functions
- Implemented in C++

- Supports:
  - Fortran77 und Fortran90, OpenMP 2.0
  - C und C++, OpenMP 1.0
  - Additional POMP directives for control and region definition
  - Used by Scalasca, VampirTrace, TAU, and ompP
  - Preserves source code information (#line line file)

- Does not support: Instrumentation of user functions

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Current Major OPARI Limitations

- Does not yet support
  - Varying number of threads in different parallel regions
  - Nested parallelism
  - Latest OpenMP 3.0 standard features like tasking

- Executed before compiler preprocessor
  - issues with macros, conditional compilation, includes!

- Needs special care if building ...
  - ... more than one application in one directory
  - ... applications spread over multiple directories

**Fixed in OPARI2**

OPARI2: will be available end of 2011

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- Measurement techniques
  - Profiling
  - Tracing

Performance Measurement

- Two dimensions
  - When performance measurement is triggered
    - External agent (asynchronous)
      - Sampling
        - Timer interrupt
        - Hardware counters overflow
        - Can measure unmodified executables, very low overhead
    - Internal agent (synchronous)
      - Code instrumentation:
        - Automatic or manual instrumentation
  - How performance data is recorded
    - Profile ::= Summation of events over time
      - run time summarization (functions, call sites, loops, …)
    - Trace file ::= Sequence of events over time
Measurement

- Typical performance data include:
  - Counts
  - Durations
  - Inclusive duration
  - Exclusive duration
  - Child duration
  - Communication cost
  - Synchronization cost
  - IO accesses
  - System calls
  - Hardware events

```
int f1()
{
    int a;
    a = a + 1;
    f2();
    a = a + 1;
    return a;
}
```

Critical Issues

- **Accuracy**
  - Perturbation
    - Measurement alters program behavior
    - E.g., memory access pattern
  - Intrusion overhead
    - Measurement itself needs time and thus lowers performance
  - Accuracy of timers, counters

- **Granularity**
  - How many measurements
  - How much information / work during each measurement

- **Tradeoff**
  - Accuracy $\leftrightarrow$ expressiveness of data
Measurement Methods: Profiling

- Recording of aggregated information
  - Time
  - Counts
    - Calls
    - Hardware counters
- about program and system entities
  - Functions, call sites, loops, basic blocks, …
  - Processes, threads
- Methods to create a profile
  - PC sampling (statistical approach)
  - Interval timer / direct measurement (deterministic approach)

Profiling (2)

- Sampling
  - General statistical measurement technique based on the assumption that a subset of a population being examined is representative for the whole population
  - Running program is interrupted periodically
    - Operating system signal or Hardware counter overflow
    - Interrupt service routine examines return-address stack to find address of instruction being executed when interrupt occurred
      - Using symbol-table information this address is mapped onto specific subroutine
  - Requires long-running programs
- Interval timing
  - Time measurement at the beginning and at the end of a code region
  - Requires instrumentation + high-resolution / low-overhead clock
**Measurement Methods: Tracing**

- Recording information about significant points (events) during execution of the program
  - Enter/leave a code region (function, loop, ...)
  - Send/receive a message...
- Save information in event record
  - Timestamp, location ID, event type
  - plus event specific information
- Event trace := stream of event records sorted by time
- Can be used to reconstruct the dynamic behavior
  ⇒ Abstract execution model on level of defined events

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**Event tracing**

**Process A**
```c
void foo() {
  trc_enter("foo");
  ...
  trc_send(B);
  send(B, tag, buf);
  ...
  trc_exit("foo");
}
```

**Process B**
```c
void bar() {
  trc_enter("bar");
  ...
  recv(A, tag, buf);
  trc_recv(A);
  ...
  trc_exit("bar");
}
```

**Local trace A**
```
58 ENTER 1
62 SEND B
64 EXIT 1
```

**Local trace B**
```
60 ENTER 1
68 RECV A
69 EXIT 1
```

**Global trace**
```
58 A ENTER 1
60 B ENTER 2
62 A SEND B
64 A EXIT 1
68 B RECV A
69 B EXIT 2
```

**Merge and unify**
```
1 foo
2 bar
```
Tracing: Time-line Visualization

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>master</td>
</tr>
<tr>
<td>2</td>
<td>slave</td>
</tr>
<tr>
<td>3</td>
<td>...</td>
</tr>
</tbody>
</table>

... 58 A ENTER 1
58 B ENTER 2
60 A SEND B
62 A EXIT 1
64 B RECV A
68 B EXIT 2
69 A SEND B
... 70 B EXIT 2

Tracing vs. Profiling

- **Tracing Advantages**
  - Event traces preserve the **temporal** and **spatial** relationships among individual events (⇒ context!)
  - Allows **reconstruction of dynamic behavior** of application on any required abstraction level
    - Automatic analysis
    - Visualization
  - Most general measurement technique
    - Profile data can be constructed from event traces
- **Disadvantages**
  - Traces can become very large
  - Writing events to a file at runtime can cause perturbation
  - Writing tracing software is complicated
    - Event buffering, clock synchronization, …
Trace File Formats

- Current Vampir trace formats
  - VTF: family of historical ASCII and binary formats
    - http://www.cs.uoregon.edu/research/paracomp/tau/vtf3-1.43.tar.gz
  - OTF: new Open Trace Format
    - http://www.tu-dresden.de/zih/otf/
- TAU performance analysis toolset
  - http://tau.uoregon.edu/docs.php#api
- EPILOG: Jülich open-source trace format
  - http://www.scalasca.org
- MPICH Multi-Processing Environment (ALOG, CLOG, SLOG, SLOG-2)
  - http://wwwunix.mcs.anl.gov/perfvis/software/log_format/
- Paraver trace analyzer (BSC, CEPBA)
  - http://www.bsc.es/paraver

No Single Solution is Sufficient!

- Combination of methods, techniques and tools needed
  - Instrumentation
    - Source code / binary, static / dynamic, manual / automatic
  - Measurement
    - Internal / external trigger, profiling / tracing
  - Analysis
    - Statistics, Visualization, Automatic, Data mining, …